

Cryptography against CONTINUAL MEMORY LEAKAGE



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Secrets

Information accessible to one party and not to other(s)
essential to cryptography!

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Theory



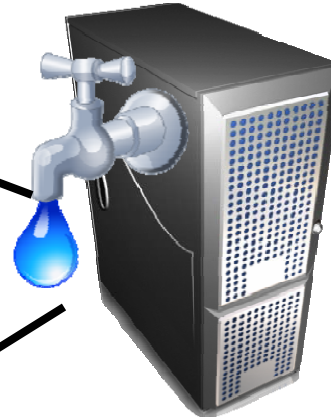
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Theory



Real life



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Theory



Real life



Secrets leak!



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- From **EVERYWHERE**
 - HD, RAM, cache, registers, randomness...
- **ALL THE TIME**
 - Not necessarily a one time process.
 - No “protected” times.



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=

Continual Memory Leakage

Previous Leakage Models

“Only Computation Leaks”

[Micali-Reyzin'04]

“Memory Leakage”

[Akavia-Goldwasser-V'09] inspired
by cold-boot attacks [HSH+08]

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Assumes a form of secure memory

Which does NOT leak as long as
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... but only ONE-SHOT

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+ Work on restricted leakage functions [Riv,CDHKS,ISW,FRRTV]

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This work:

**remove these
restrictions!**

Our Model: Continual Memory Leakage



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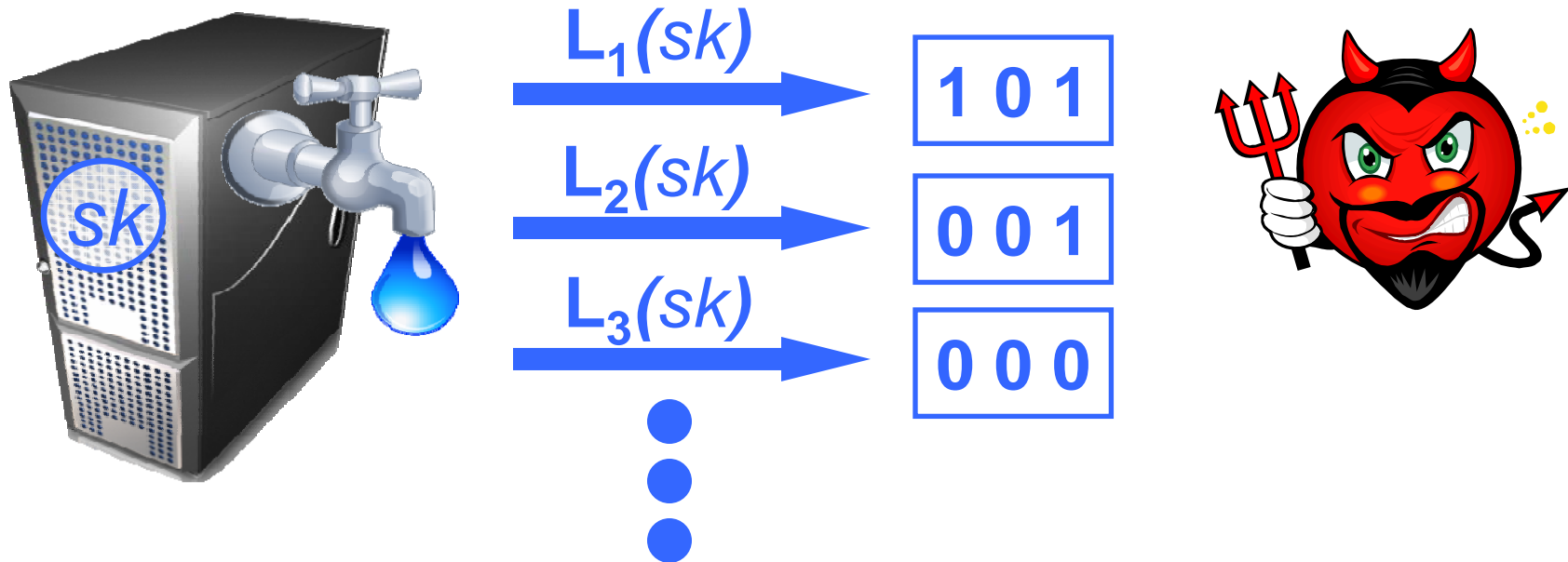


$L_1(sk)$

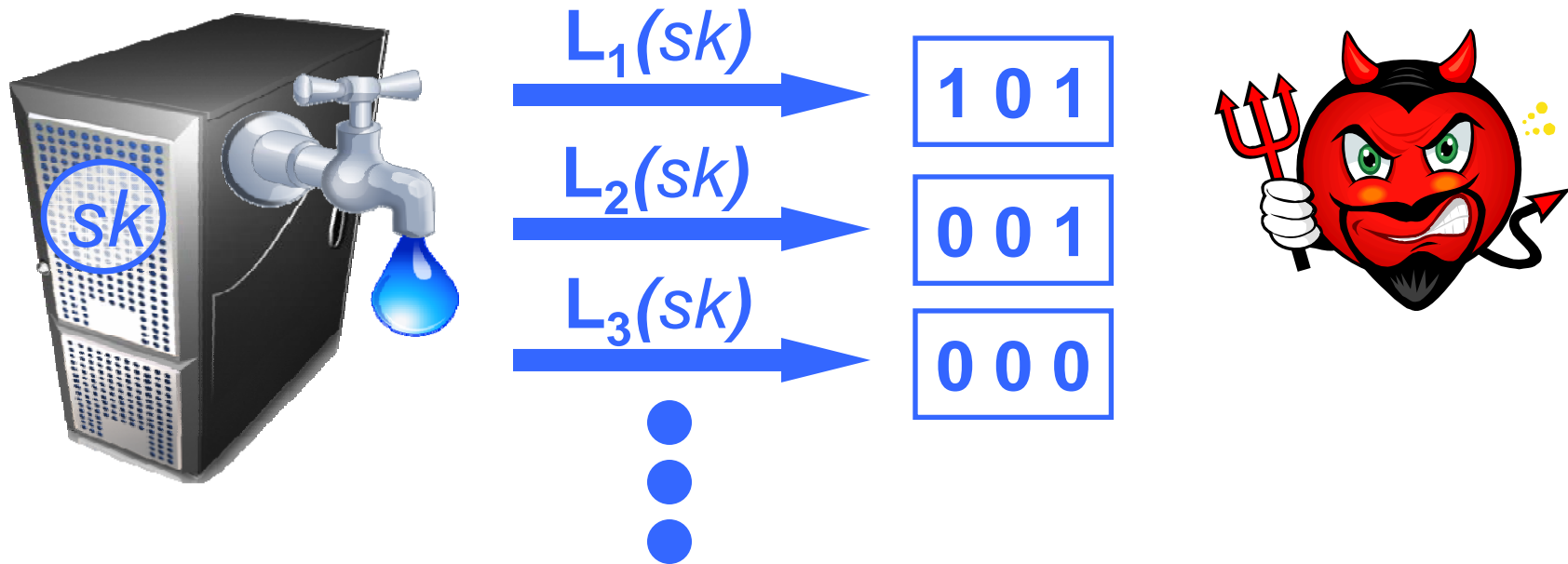
1 0 1



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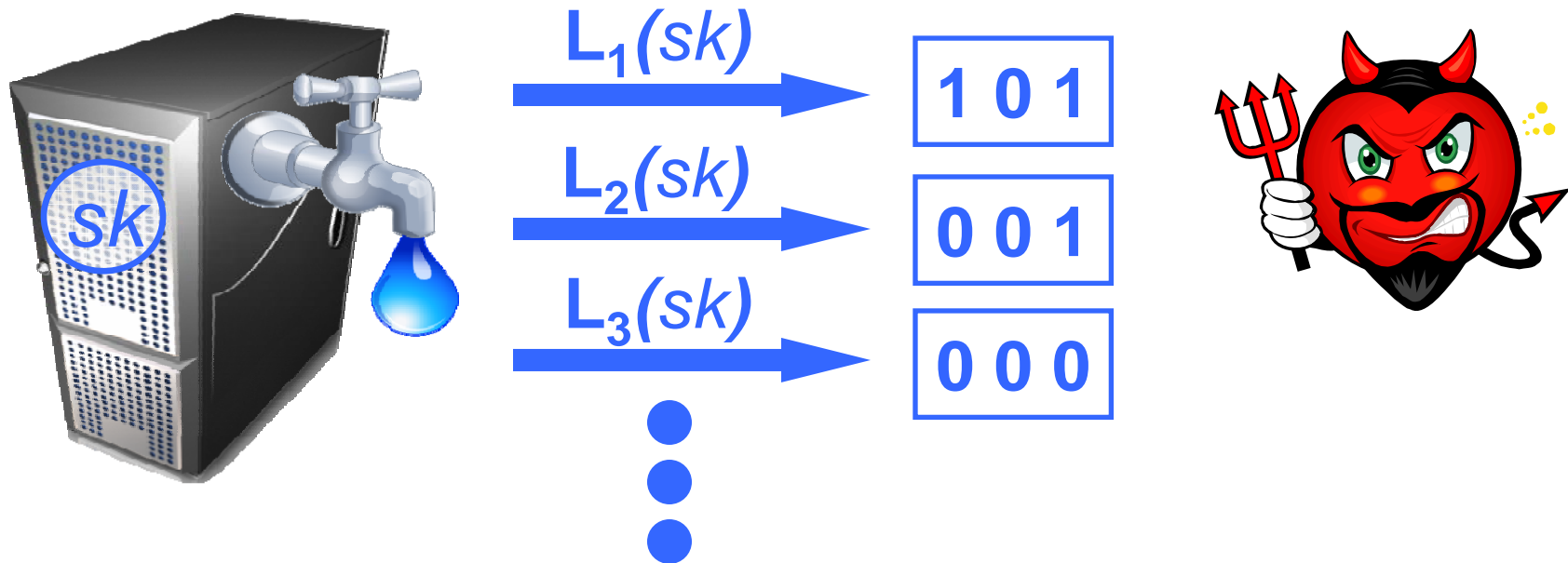


Our Model: Continual Memory Leakage



Can we protect against continual leakage?

Our Model: Continual Memory Leakage



Can we protect against continual leakage?

Need to (periodically) update SK .

Our Model: Continual Memory Leakage



$L_1(sk_1)$



1 0 1



Our Model: Continual Memory Leakage



$L_1(sk_1)$

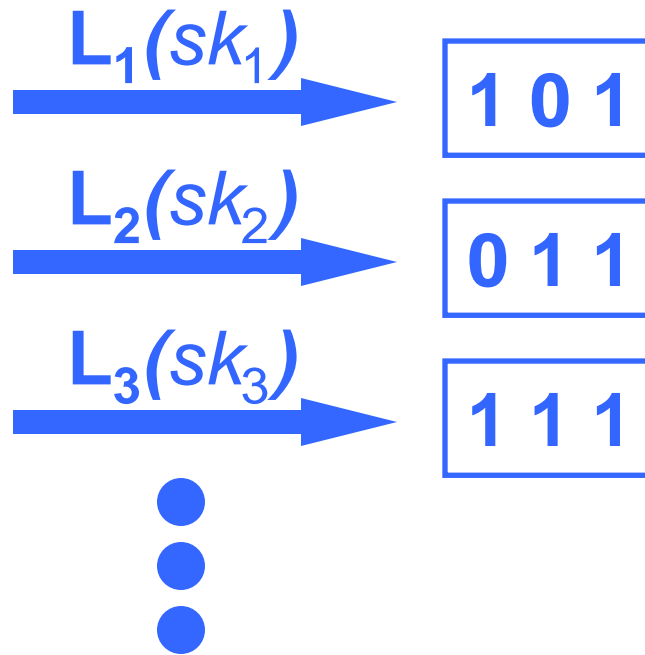
$L_2(sk_2)$

1 0 1

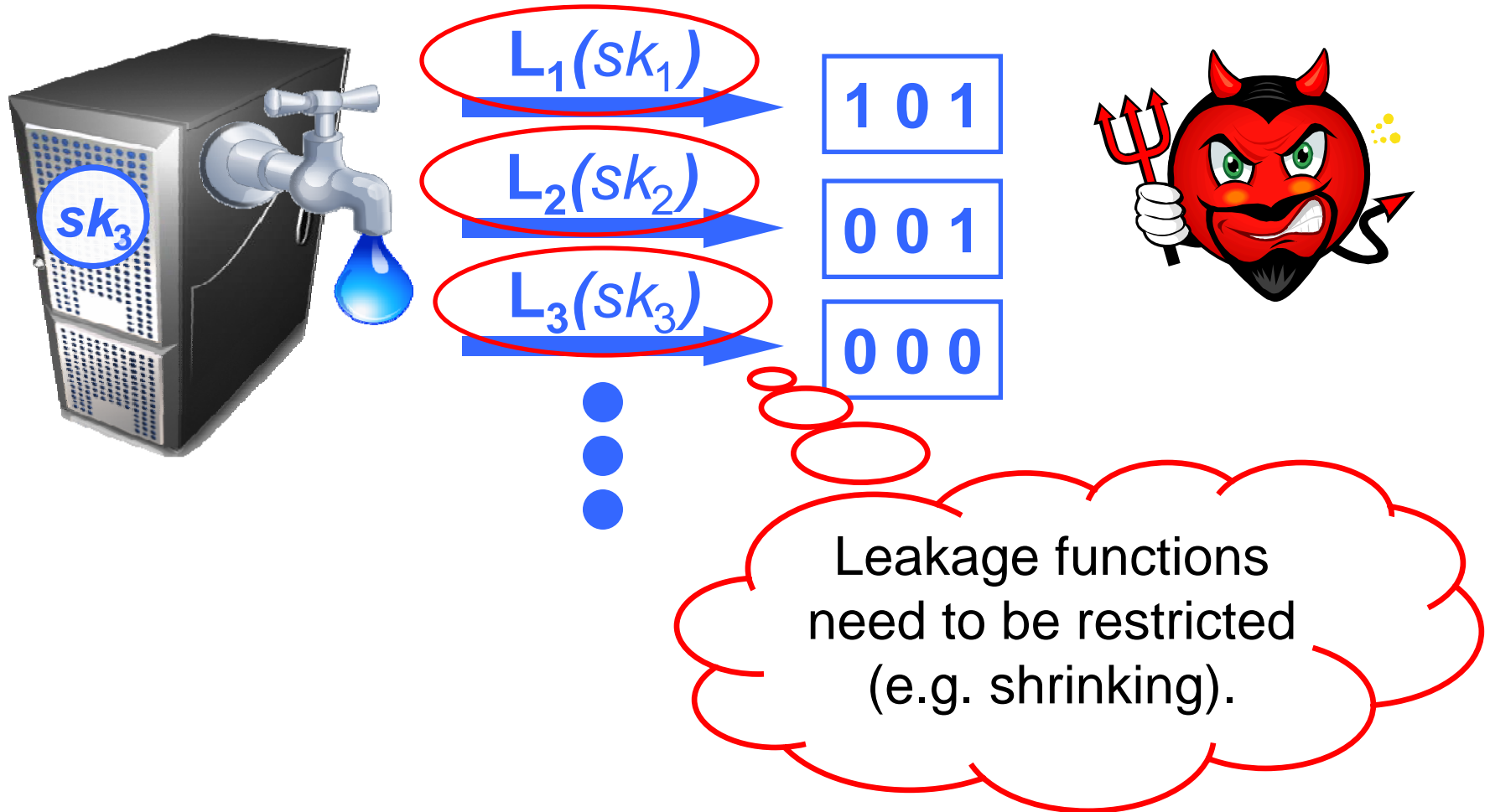
0 1 1



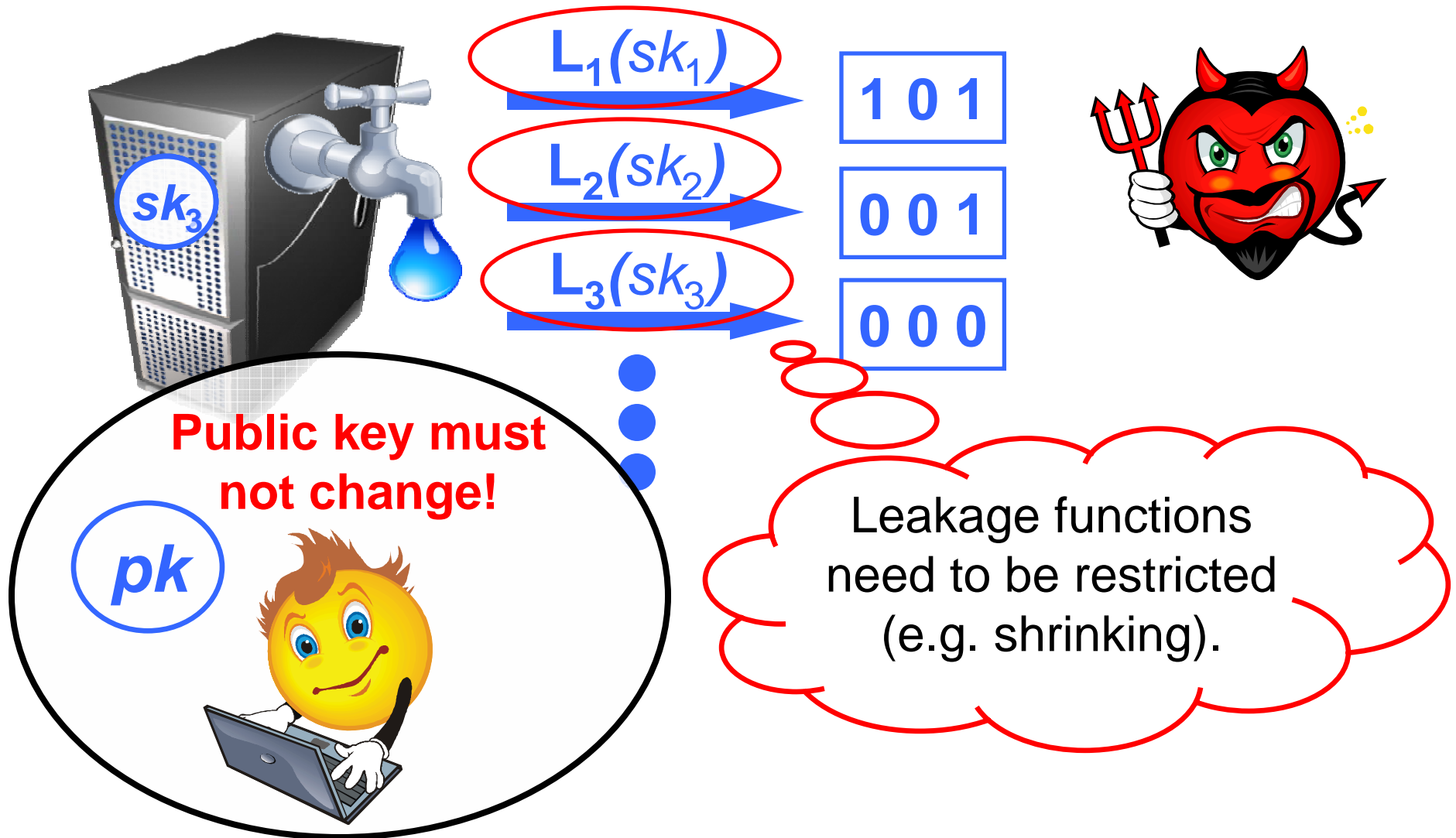
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Our Model: Continual Memory Leakage



Features of Continual Memory Leakage (CML)

- Leakage function applied to **entire** memory
- Key updates are oblivious to users (public key doesn't change).
- Leakage can occur at any time point!
 - Including during **key updates**.
 - Including during **decrypting/signing**.
- Total amount of leakage – **unbounded**.
 - Only the rate (leaked bits/sec) is bounded.



Our Results: Cryptography in the CML Model

- **Public-key encryption.**
 - **Linear assumption:** leakage rate $(1/2-o(1))$.
 - **SXDH assumption:** leakage rate $(1-o(1))$.
 - **KeyGen and Update:** logarithmic no. of bits.
- Identity based encryption.
- Encryption \Rightarrow Signatures.

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Previously not known even
in the Micali-Reyzin model!

(under standard assumptions)

- Identical
- Encryption \Rightarrow Signature

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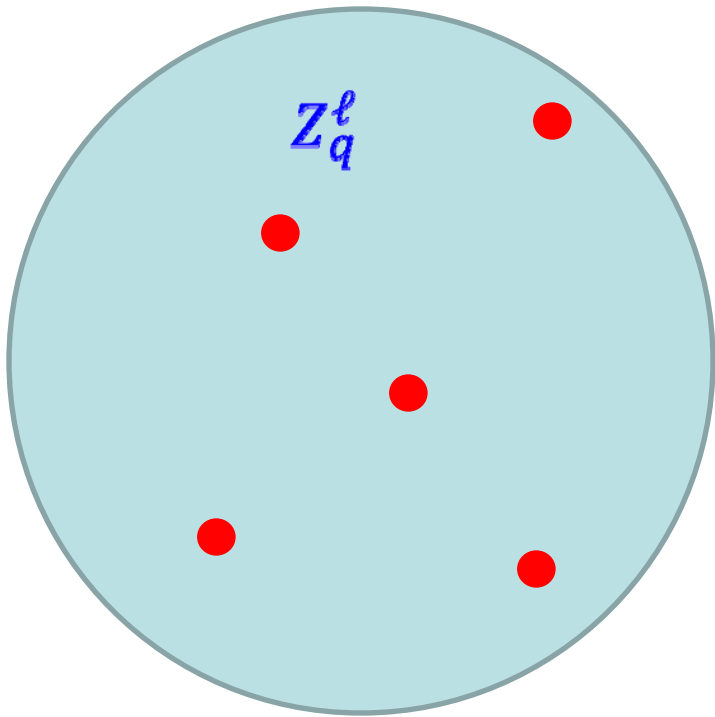
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Concurrently [DHLW]: *efficient* signatures, ID schemes and AKA in the CML model under linear and SXDH.

- Different techniques (re-randomizable NIZK).

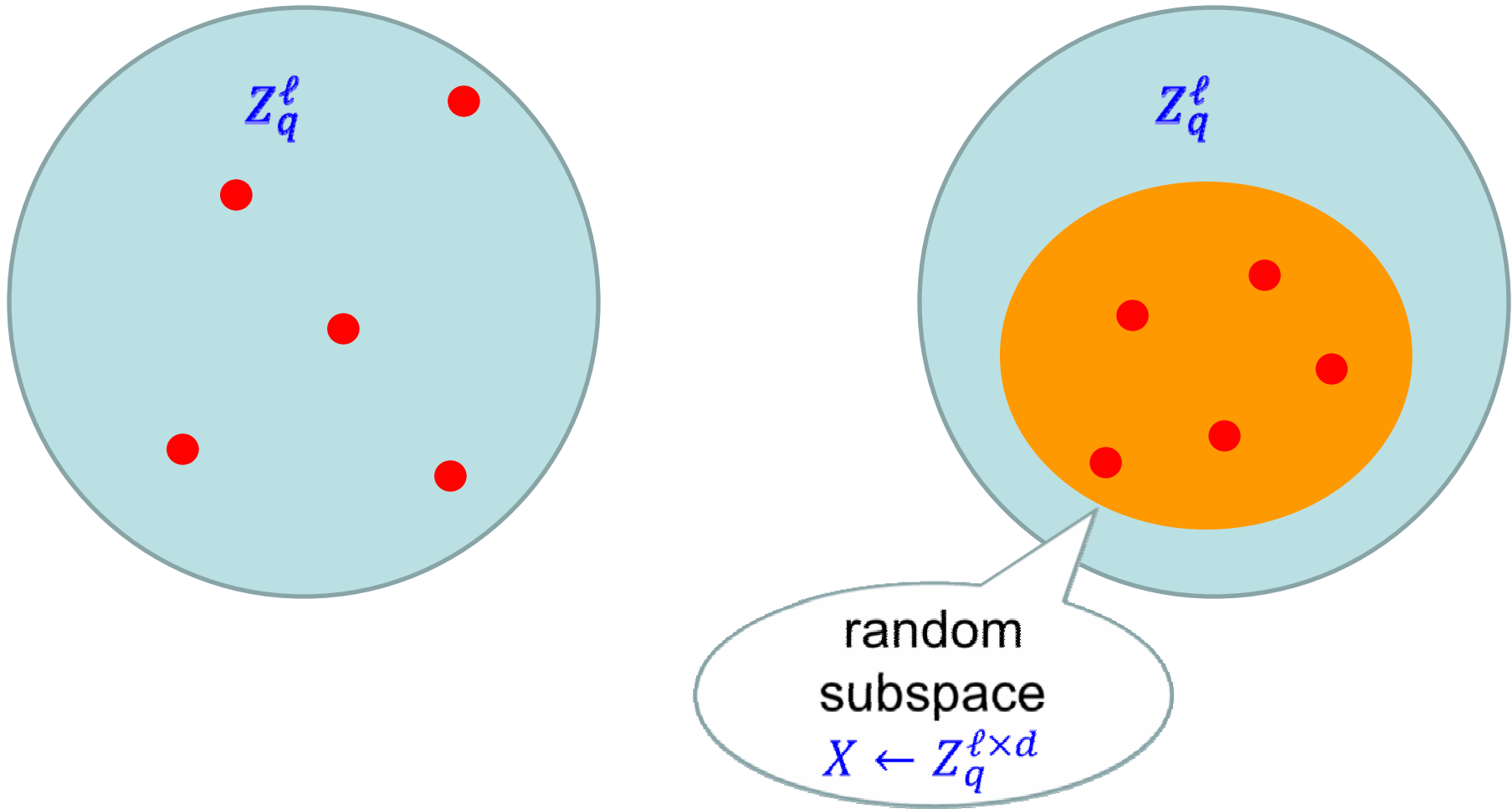
A Sneak Peek

Lemma: Random Subspaces are Leakage-Resilient



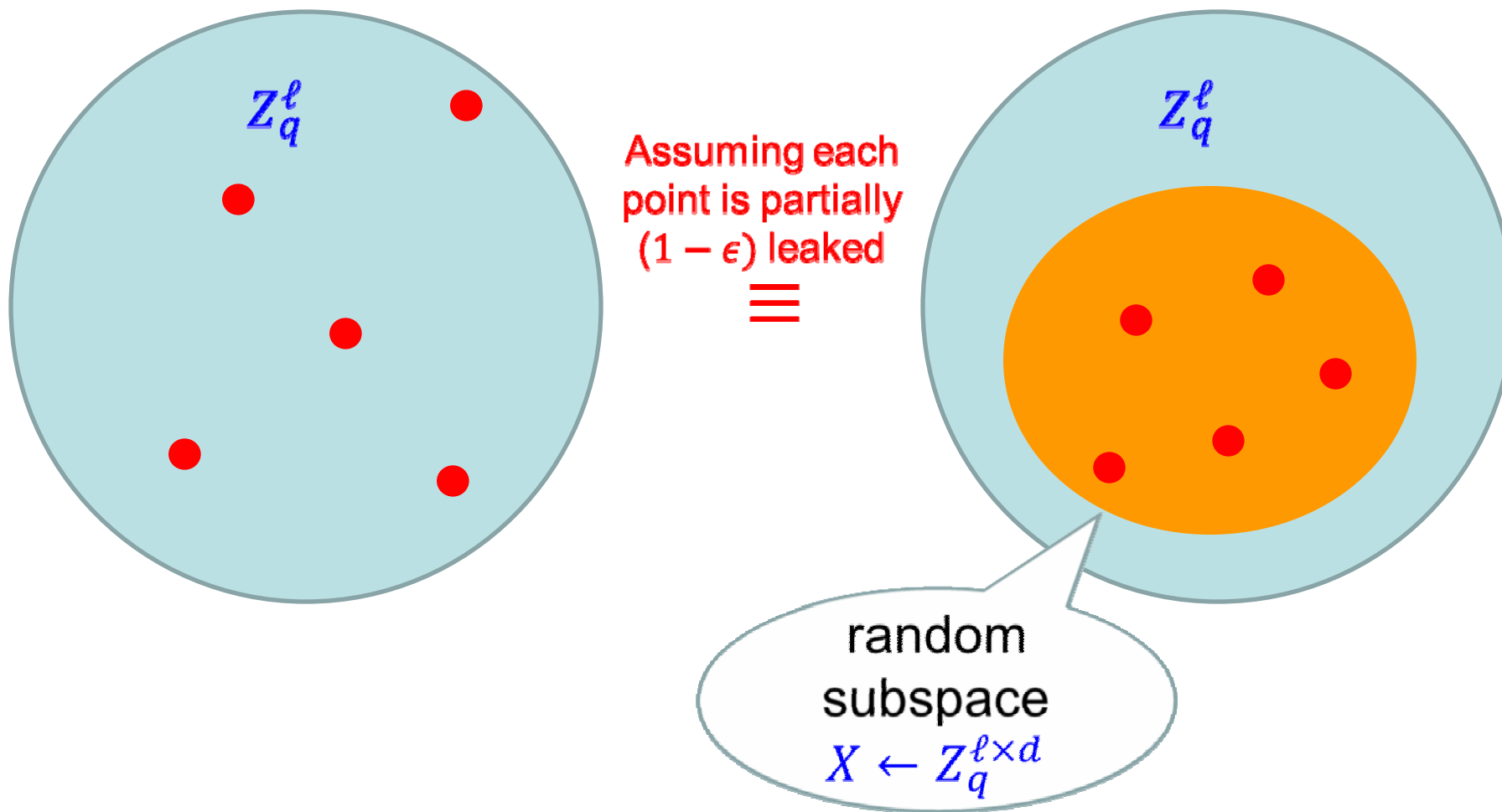
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Paper on Eprint:

“Cryptography against Continual Memory Leakage”,

<http://eprint.iacr.org/2010/278>.

